

*Project:* DRA FRONT END FILTER PROJECT

*Title:* Proof of Security (IIa)

*Ref:* DS/FMU/FEF/010

*Issue: Revision : 2.4*

*Date:* 5 June 2016

*Status:* Approved

*Type:* Specification

*Keywords:*

*Author:*

<i>Name</i>	<i>Location</i>	<i>Signature</i>	<i>Date</i>
G. M. Prout	WIN01		

*Authorisation for Issue:*

<i>Name</i>	<i>Function</i>	<i>Signature</i>	<i>Date</i>
R.B. Jones	HAT Manager		

*Abstract:* This document provides a formal proof for the security property on *hide* for the DRA front end filter project RSRE 1C/6130.

*Distribution:* HAT FEF File  
Simon Wiseman

---

## 0 DOCUMENT CONTROL

### 0.1 Contents List

<b>0</b>	<b>DOCUMENT CONTROL</b>	<b>2</b>
0.1	Contents List . . . . .	2
0.2	Document Cross References . . . . .	3
0.3	Changes History . . . . .	3
0.4	Changes Forecast . . . . .	3
<b>1</b>	<b>GENERAL</b>	<b>4</b>
1.1	Scope . . . . .	4
1.2	Introduction . . . . .	4
<b>2</b>	<b>PRELIMINARIES</b>	<b>4</b>
<b>3</b>	<b>SETTING UP FOR THE SECURITY PROOF</b>	<b>4</b>
3.1	Theorem about <i>dominates</i> . . . . .	4
3.2	Consistency Proof for <i>repState</i> and <i>absState</i> . . . . .	5
3.3	Retrieving the Remaining Definitions of Constants . . . . .	5
3.4	Components of Labelled Products . . . . .	6
3.4.1	Type <i>Directory</i> . . . . .	6
3.4.2	Type <i>TableSpec</i> . . . . .	7
3.4.3	Type <i>Row</i> . . . . .	8
3.4.4	Type <i>Data</i> . . . . .	8
<b>4</b>	<b>PROOF OF SECURITY OF <i>hide</i></b>	<b>9</b>
4.1	Proof Strategy . . . . .	9
4.2	Proof of Security of <i>hideR</i> . . . . .	10
4.2.1	<i>replaceData</i> Lemma . . . . .	11
4.2.2	<i>cleanColCons</i> Lemma . . . . .	12
4.2.3	<i>cleanRow</i> Lemma . . . . .	13
4.2.4	<i>cleanTable</i> Lemma . . . . .	16
4.2.5	<i>cleanDirectory</i> Lemma . . . . .	22
4.2.6	<i>hideR</i> Lemma . . . . .	24
4.3	Proof that the Invariant on the State is Maintained . . . . .	26
4.4	<i>hide</i> Lemma . . . . .	30
<b>5</b>	<b>CLOSING DOWN</b>	<b>31</b>
<b>6</b>	<b>THE THEORY fef010</b>	<b>32</b>
6.1	Parents . . . . .	32
6.2	Children . . . . .	32
6.3	Constants . . . . .	32
6.4	Definitions . . . . .	32
6.5	Theorems . . . . .	32
<b>7</b>	<b>INDEX</b>	<b>36</b>

## 0.2 Document Cross References

- [1] DS/FMU/017. *Secure Database Technical Proposal*. High Assurance Team, ICL Secure Systems, WIN01, 21st January 1992.
- [2] DS/FMU/FEF/007. *Proof Strategy*. G.M. Prout, ICL Secure Systems, WIN01.

## 0.3 Changes History

**Issue Revision : 2.4 (5 June 2016)** Consistency proof for *dominates* extended because of inclusion of *glb*.

**Issue 2.5** Removed dependency on ICL logo font

## 0.4 Changes Forecast

None.

## 1 GENERAL

### 1.1 Scope

This document provides a formal proof that the component *hide* satisfies its critical requirement, as specified in the proof strategy [2]. It constitutes part of deliverable D6 of work package 1c, as given in section 7 of the Secure Database Technical Proposal, [1].

### 1.2 Introduction

This document is a proof script which provides a formal proof of the first conjunct of *Lemma1*, the requirement on the critical components *hide* and *updateState*, described in the proof strategy document [2].

#### Lemma1

$$\vdash \quad \text{hide} \in \text{secureHide} \wedge (\text{hide}, \text{updateState}) \in \text{secureUpdate}$$

In this document, we give a proof of

$$\vdash \quad \text{hide} \in \text{secureHide}$$

First, we define a property *secureHideR* on the representation type, and then prove

$$\vdash \quad \text{hideR} \in \text{secureHideR}$$

We then prove that *hideR* maintains the state invariant. From these two results, we prove that *hide* is secure.

## 2 PRELIMINARIES

The following ProofPower instructions set up the new theory *fef010*.

SML

```
| open_theory "fef009";
| (force_delete_theory "fef010" handle _ => ());
| new_theory "fef010";
| push_merge_pcs["hol","wrk049","pair1"] ;
```

## 3 SETTING UP FOR THE SECURITY PROOF

### 3.1 Theorem about *dominates*

A theorem which just states the transitivity property of *dominates*.

SML

```
| val dominates_trans = prove_rule[dominates_def]
|    $\ulcorner \forall x y z \bullet x \text{ dominates } y \wedge y \text{ dominates } z \Rightarrow x \text{ dominates } z \urcorner$ ;
```

HOL output

```
| dominates_trans =  $\vdash \forall x y z \bullet x \text{ dominates } y \wedge y \text{ dominates } z \Rightarrow x \text{ dominates } z$ 
```

### 3.2 Consistency Proof for *repState* and *absState*

We prove the consistency of *repState* and *absState* and retrieve their definitions with the consistency obligation satisfied.

SML

```
| push_consistency_goal  $\ulcorner \text{repState} \urcorner$ ;
| a(strip_asm_tac (simple_ $\Rightarrow$ _match_mp_rule type_lemmas_thm state_type_def));
| a( $\exists$ _tac  $\ulcorner (\text{rep}, \text{abs}) \urcorner$  THEN asm_rewrite_tac[]);
| a  $\wedge$ _tac;
```

SML

```
| (* *** Goal "1" *** *)
| a(REPEAT strip_tac THEN_TRY asm_rewrite_tac[]);
| a(POP_ASM_T (ante_tac o app_fun_rule  $\ulcorner \text{abs} \urcorner$ ) THEN asm_rewrite_tac[]);
| (* *** Goal "2" *** *)
| a(REPEAT  $\forall$ _tac THEN  $\Rightarrow$ _tac);
| a(REPEAT strip_tac THEN_TRY asm_rewrite_tac[]);
| a(POP_ASM_T (ante_tac o app_fun_rule  $\ulcorner \text{rep} \urcorner$ ) THEN asm_rewrite_tac[]);
| save_consistency_thm  $\ulcorner \text{repState} \urcorner$  (pop_thm());
| val repState_absState_def = get_spec  $\ulcorner \text{repState} \urcorner$ ;
```

### 3.3 Retrieving the Remaining Definitions of Constants

SML

```
| val hide_def = get_spec  $\ulcorner \text{hide} \urcorner$ ;
| val hideR_def = get_spec  $\ulcorner \text{hideR} \urcorner$ ;
| val secureHide_def = get_spec  $\ulcorner \text{secureHide} \urcorner$ ;
| val cleanDirectory_def = get_spec  $\ulcorner \text{cleanDirectory} \urcorner$ ;
| val cleanTable_def = conv_rule(MAP_C let_conv)(get_spec  $\ulcorner \text{cleanTable} \urcorner$ );
| val cleanColCons_def = conv_rule(MAP_C let_conv)(get_spec  $\ulcorner \text{cleanColCons} \urcorner$ );
| val cleanRows_def = get_spec  $\ulcorner \text{cleanRows} \urcorner$ ;
| val cleanRow_def = get_spec  $\ulcorner \text{cleanRow} \urcorner$ ;
| val replaceData_def = get_spec  $\ulcorner \text{replaceData} \urcorner$ ;
| val filterRow_def = get_spec  $\ulcorner \text{filterRow} \urcorner$ ;
```

### 3.4 Components of Labelled Products

These theorems state that if two things of labelled product type are equal, then their components must be equal.

#### 3.4.1 Type *Directory*

SML

```

push_goal([],Γ∀ dir1 dir2 •
  dir1 = dir2 ⇔
  Dir_tables dir1 = Dir_tables dir2
  ∧ Dir_exist dir1 = Dir_exist dir2
  ∧ Dir_class dir1 = Dir_class dir2⌈);
a(REPEAT strip_tac THEN_TRY asm_rewrite_tac[]);
a(lemma_tacΓ(dir1 = MkDirectory (Dir_tables dir1) (Dir_exist dir1) (Dir_class dir1))
  ∧ (dir2 = MkDirectory (Dir_tables dir2) (Dir_exist dir2) (Dir_class dir2))⌈);
(* *** Goal "1" *** *)
a(rewrite_tac[get_specΓMkDirectory⌈]);
(* *** Goal "2" *** *)
a(POP_ASM_T once_rewrite_thm_tac);
a(POP_ASM_T once_rewrite_thm_tac);
a(asm_rewrite_tac[]);
val dir_components = save_pop_thm"dir_components";

```

HOL output

```

dir_components =
⊢ ∀ dir1 dir2
  • dir1 = dir2
    ⇔ Dir_tables dir1 = Dir_tables dir2
      ∧ Dir_exist dir1 = Dir_exist dir2
      ∧ Dir_class dir1 = Dir_class dir2

```

**3.4.2 Type TableSpec**

SML

```

push_goal([],Γ∀ t1 t2 •
  t1 = t2 ⇔
  TS_class t1 = TS_class t2
∧ TS_maxRow t1 = TS_maxRow t2
∧ TS_colspecs t1 = TS_colspecs t2
∧ TS_cons t1 = TS_cons t2
∧ TS_rows t1 = TS_rows t2¬);
a(REPEAT strip_tac THEN_TRY asm_rewrite_tac[]);
a(lemma_tacΓ(t1 = MkTableSpec(TS_class t1) (TS_maxRow t1) (TS_colspecs t1)
  (TS_cons t1)(TS_rows t1))
  ∧ (t2 = MkTableSpec(TS_class t2) (TS_maxRow t2) (TS_colspecs t2)
  (TS_cons t2)(TS_rows t2))¬);
(* *** Goal "1" *** *)
a(rewrite_tac[get_specΓMkTableSpec¬]);
(* *** Goal "2" *** *)
a(POP_ASM_T once_rewrite_thm_tac);
a(POP_ASM_T once_rewrite_thm_tac);
a(asm_rewrite_tac[]);
val tab_components = save_pop_thm"tab_components";

```

HOL output

```

tab_components =
| ⊢ ∀ t1 t2
| • t1 = t2
|   ⇔ TS_class t1 = TS_class t2
|     ∧ TS_maxRow t1 = TS_maxRow t2
|     ∧ TS_colspecs t1 = TS_colspecs t2
|     ∧ TS_cons t1 = TS_cons t2
|     ∧ TS_rows t1 = TS_rows t2

```

### 3.4.3 Type Row

SML

```

push_goal([],Γ∀ r1 r2 •
  r1 = r2 ⇔
  R_exist r1 = R_exist r2
  ∧ R_data r1 = R_data r2⌈);
a(REPEAT strip_tac THEN_TRY asm_rewrite_tac[]);
a(lemma_tacΓ(r1 = MkRow(R_exist r1) (R_data r1))
  ∧ (r2 = MkRow(R_exist r2) (R_data r2))⌈);
(* *** Goal "1" *** *)
a(rewrite_tac[get_specΓMkRow⌈]);
(* *** Goal "2" *** *)
a(POP_ASM_T once_rewrite_thm_tac);
a(POP_ASM_T once_rewrite_thm_tac);
a(asm_rewrite_tac[]);
val row_components = save_pop_thm"row_components";

```

HOL output

```

row_components =
  ⊢ ∀ r1 r2
  • r1 = r2 ⇔ R_exist r1 = R_exist r2 ∧ R_data r1 = R_data r2

```

### 3.4.4 Type Data

SML

```

push_goal([],Γ∀ d1 d2 •
  d1 = d2 ⇔
  Dat_class d1 = Dat_class d2
  ∧ Dat_item d1 = Dat_item d2⌈);
a(REPEAT strip_tac THEN_TRY asm_rewrite_tac[]);
a(lemma_tacΓ(d1 = MkData(Dat_class d1) (Dat_item d1))
  ∧ (d2 = MkData(Dat_class d2) (Dat_item d2))⌈);
(* *** Goal "1" *** *)
a(rewrite_tac[get_specΓMkData⌈]);
(* *** Goal "2" *** *)
a(POP_ASM_T once_rewrite_thm_tac);
a(POP_ASM_T once_rewrite_thm_tac);
a(asm_rewrite_tac[]);
val data_components = save_pop_thm"data_components";

```

HOL output

```

val data_components =
  ⊢ ∀ d1 d2
  • d1 = d2 ⇔ Dat_class d1 = Dat_class d2 ∧ Dat_item d1 = Dat_item d2

```



---

## 4 PROOF OF SECURITY OF *hide*

### 4.1 Proof Strategy

First we define *secureHideR* a property on things of type  $Class \times StateR \rightarrow StateR$

HOL Constant

$\mathbf{secureHideR} : (Class \times StateR \rightarrow StateR) \mathbb{P}$ <hr style="border: 0.5px solid black; margin: 10px 0;"/> $\forall h : Class \times StateR \rightarrow StateR \bullet$ $h \in \mathit{secureHideR}$ $\Leftrightarrow$ $\forall c_1 c_2 : Class; s_1 s_2 : StateR \bullet$ $h(c_1, s_1) = h(c_1, s_2)$ $\wedge$ $c_1 \text{ dominates } c_2$ $\Rightarrow$ $h(c_2, s_1) = h(c_2, s_2)$
---

We simplify some of the properties of the representation and abstraction functions.

SML

```
| val isState_lemma = all_∀_intro(nth 4 (strip_∧_rule(all_∀_elim repState_absState_def)));
```

HOL output

```
| isState_lemma = ⊢ ∀ s • isState (repState s)
```

SML

```
| val isState_lemma1 = all_∀_intro(nth 1 (strip_∧_rule(all_∀_elim repState_absState_def)));
```

HOL output

```
| isState_lemma1 = ⊢ isState r ⇒ repState (absState r) = r
```

SML

```
| val isState_lemma2 = all_∀_intro(nth 3 (strip_∧_rule(all_∀_elim repState_absState_def)));
```

HOL output

```
| isState_lemma2 =
| ⊢ ∀ r1 r2
| • isState r1 ∧ isState r2 ⇒ (absState r1 = absState r2 ⇔ r1 = r2)
```

Next, we prove that if *hideR* is a member of *secureHideR* and *hideR* maintains the invariant on the representation state, then *hide* is a member of *secureHide*.

SML

```

push_goal([],⌈(hideR ∈ secureHideR ∧ ∀ clear s • isState s ⇒ isState(hideR(clear,s)))
  ⇒ hide ∈ secureHide⌋);
a(rewrite_tac[secureHide_def,hide_def,get_spec⌈secureHideR⌋] THEN REPEAT strip_tac);
a(DROP_NTH_ASM_T 4 (ante_tac o list_∀_elim[⌈c1⌋,⌈c2⌋,⌈repState s1⌋,⌈repState s2⌋]));
a(asm_rewrite_tac[] THEN ⇒_T asm_tac);
a(GET_NTH_ASM_T 4 (asm_tac o rewrite_rule[isState_lemma] o
  list_∀_elim[⌈c1⌋,⌈repState s1⌋]));
a(GET_NTH_ASM_T 5 (asm_tac o rewrite_rule[isState_lemma] o
  list_∀_elim[⌈c1⌋,⌈repState s2⌋]));
a(strip_asm_tac (list_∀_elim[⌈hideR (c1, repState s1)⌋,⌈hideR (c1, repState s2)⌋]
  isState_lemma2));
a(DROP_NTH_ASM_T 4 ante_tac THEN asm_rewrite_tac[] THEN ⇒_tac);
a(GET_NTH_ASM_T 7 (asm_tac o rewrite_rule[isState_lemma] o
  list_∀_elim[⌈c2⌋,⌈repState s1⌋]));
a(GET_NTH_ASM_T 8 (asm_tac o rewrite_rule[isState_lemma] o
  list_∀_elim[⌈c2⌋,⌈repState s2⌋]));
a(strip_asm_tac (list_∀_elim[⌈hideR (c2, repState s1)⌋,⌈hideR (c2, repState s2)⌋]
  isState_lemma2));
val hideR_hide_lemma = save_pop_thm"hideR_hide_lemma";

```

HOL output

```

hideR_hide_lemma =
⊢ hideR ∈ secureHideR
  ∧ (∀ clear s • isState s ⇒ isState (hideR (clear, s)))
  ⇒ hide ∈ secureHide

```

## 4.2 Proof of Security of *hideR*

The proof is broken down into a series of lemmas.

**4.2.1** *replaceData* Lemma

SML

```

push_goal([],Γ∀ c1 c2 • c1 dominates c2
  ⇒ ∀ d1 d2 •
    replaceData c1 d1 = replaceData c1 d2
    ⇒
    replaceData c2 d1 = replaceData c2 d2∇);
a(REPEAT strip_tac);
a(POP_ASM_T ante_tac THEN rewrite_tac[replaceData_def]);
a(cases_tacΓc2 dominates Dat_class d1∇ THEN cases_tacΓc2 dominates Dat_class d2∇
  THEN asm_rewrite_tac[]);
(* *** Goal "1" *** *)
a(lemma_tacΓ(c1 dominates Dat_class d1) ∧ (c1 dominates Dat_class d2)∇
  THEN_LIST[basic_res_tac2 3[dominates_trans],asm_rewrite_tac[]]);

```

SML

```

(* *** Goal "2" *** *)
a(lemma_tacΓc1 dominates Dat_class d1∇ THEN_LIST
  [contr_tac THEN basic_res_tac4 2 [1][3,4][][dominates_trans],asm_rewrite_tac[]]);
a(cases_tacΓc1 dominates Dat_class d2∇ THEN asm_rewrite_tac[]);
a ⇒ _tac;
a(DROP_NTH_ASM_T 5 ante_tac THEN asm_rewrite_tac[]);
(* *** Goal "3" *** *)
a(lemma_tacΓc1 dominates Dat_class d2∇ THEN_LIST
  [contr_tac THEN basic_res_tac4 2 [1][2,4][][dominates_trans],asm_rewrite_tac[]]);
a(cases_tacΓc1 dominates Dat_class d1∇ THEN asm_rewrite_tac[]);
a ⇒ _tac;
a(DROP_NTH_ASM_T 5 ante_tac THEN asm_rewrite_tac[]);

```

SML

```

(* *** Goal "4" *** *)
a(cases_tacΓc1 dominates Dat_class d1∇ THEN cases_tacΓc1 dominates Dat_class d2∇
  THEN asm_rewrite_tac[]);
(* *** Goal "4.1" *** *)
a ⇒ _tac;
a(DROP_NTH_ASM_T 5 ante_tac THEN asm_rewrite_tac[]);
(* *** Goal "4.2" *** *)
a(⇒_T(strip_asm_tac o rewrite_rule[data_components,get_specΓMkData∇])); ;
a(asm_rewrite_tac[]);
(* *** Goal "4.3" *** *)
a(⇒_T(strip_asm_tac o rewrite_rule[data_components,get_specΓMkData∇])); ;
a(asm_rewrite_tac[]);
val replaceData_lemma = save_pop_thm"replaceData_lemma";

```

HOL output

```

replaceData_lemma =
| ⊢ ∀ c1 c2
|   • c1 dominates c2
|     ⇒ (∀ d1 d2
|       • replaceData c1 d1 = replaceData c1 d2
|         ⇒ replaceData c2 d1 = replaceData c2 d2)

```

**4.2.2 cleanColCons Lemma**

SML

```

push_goal([], ⊢ ∀ c1 c2 t1 t2 •
|   (c1 dominates c2
|     ∧ cleanColCons c1 t1 = cleanColCons c1 t2)
|     ⇒ cleanColCons c2 t1 = cleanColCons c2 t2);
a(REPEAT strip_tac);
a(POP_ASM_T ante_tac THEN rewrite_tac[cleanColCons_def, ▷_def, dom_def]);
a(PC_T "hol1"(REPEAT strip_tac));
(* *** Goal "1" *** *)
a(lemma_tac ⊢ c1 dominates CC_exist (Snd x));
(* *** Goal "1.1" *** *)
a(ante_tac(list_∀_elim[⊢ c1, ⊢ c2, ⊢ CC_exist (Snd x)]dominates_def) THEN asm_rewrite_tac[]);
a(REPEAT strip_tac);
(* *** Goal "1.2" *** *)
a(PC_T"hol1"(DROP_NTH_ASM_T 5 (strip_asm_tac o ∀_elim ⊢ x)));

```

SML

```

(* *** Goal "2" *** *)
a(lemma_tac ⊢ c1 dominates CC_exist (Snd x));
(* *** Goal "2.1" *** *)
a(ante_tac(list_∀_elim[⊢ c1, ⊢ c2, ⊢ CC_exist (Snd x)]dominates_def) THEN asm_rewrite_tac[]);
a(REPEAT strip_tac);
(* *** Goal "2.2" *** *)
a(PC_T"hol1"(DROP_NTH_ASM_T 5 (strip_asm_tac o ∀_elim ⊢ x)));
(* *** Goal "3" *** *)
a(∃_tac ⊢ y THEN asm_rewrite_tac[]);
a(lemma_tac ⊢ c1 dominates CC_exist y);
(* *** Goal "3.1" *** *)
a(ante_tac(list_∀_elim[⊢ c1, ⊢ c2, ⊢ CC_exist y]dominates_def) THEN asm_rewrite_tac[]);
a(REPEAT strip_tac);
(* *** Goal "3.2" *** *)
a(PC_T"hol1"(DROP_NTH_ASM_T 5 (strip_asm_tac o ∀_elim ⊢ (CS_consGroup x, y))));

```

SML

```

(* *** Goal "4" *** *)
a( $\exists$ -tac $\ulcorner$ y $\urcorner$  THEN asm_rewrite_tac[]);
a(lemma_tac  $\ulcorner$ c1 dominates CC_exist y $\urcorner$ );
(* *** Goal "4.1" *** *)
a(ante_tac(list_ $\forall$ -elim $\ulcorner$ c1 $\urcorner$ , $\ulcorner$ c2 $\urcorner$ , $\ulcorner$ CC_exist y $\urcorner$ )dominates_def) THEN asm_rewrite_tac[]);
a(REPEAT strip_tac);
(* *** Goal "4.2" *** *)
a(PC-T"hol1"(DROP-NTH-ASM-T 5 (strip_asm_tac o  $\forall$ -elim $\ulcorner$ (CS_consGroup x, y) $\urcorner$ )));
val cleanColCons_lemma = save_pop_thm"cleanColCons_lemma";

```

HOL output

```

cleanColCons_lemma =
 $\vdash \forall c_1 c_2 t_1 t_2$ 
  •  $c_1$  dominates  $c_2 \wedge$  cleanColCons  $c_1 t_1 =$  cleanColCons  $c_1 t_2$ 
     $\Rightarrow$  cleanColCons  $c_2 t_1 =$  cleanColCons  $c_2 t_2$ 

```

### 4.2.3 cleanRow Lemma

First, two results about *cleanColCons*.

SML

```

push_goal([], $\ulcorner$  $\forall c_1 c_2 t$  •  $c_1$  dominates  $c_2$ 
   $\Rightarrow$  ( $Fst(cleanColCons c_2 t) \subseteq Fst(cleanColCons c_1 t)$ 
     $\wedge Snd(cleanColCons c_2 t) \subseteq Snd(cleanColCons c_1 t)$ ) $\urcorner$ );
a(rewrite_tac[rel_ext_clauses,sets_ext_clauses,cleanColCons_def, $\triangleright$ -thm,dom_def]
  THEN REPEAT strip_tac);
(* *** Goal "1" *** *)
a(contr_tac THEN basic_res_tac4 2 [1][3,4][][][dominates_trans]);
(* *** Goal "2" *** *)
a( $\exists$ -tac $\ulcorner$ y $\urcorner$  THEN asm_rewrite_tac[]);
a(contr_tac THEN basic_res_tac4 2 [1][3,4][][][dominates_trans]);
val  $\subseteq$ -cleanColCons_lemma = save_pop_thm" $\subseteq$ -cleanColCons_lemma";

```

HOL output

```

 $\subseteq$ -cleanColCons_lemma =
 $\vdash \forall c_1 c_2 t$ 
  •  $c_1$  dominates  $c_2$ 
     $\Rightarrow Fst(cleanColCons c_2 t) \subseteq Fst(cleanColCons c_1 t)$ 
       $\wedge Snd(cleanColCons c_2 t) \subseteq Snd(cleanColCons c_1 t)$ 

```

SML

```

push_goal([],Γ∀ c1 c2 t col • c1 dominates c2
  ⇒ (col ∈ Snd(cleanColCons c2 t) ⇒ col ∈ Snd(cleanColCons c1 t))∇);
a(REPEAT strip_tac);
a(basic_res_tac 2 [rewrite_rule[sets_ext_clauses] ⊆_cleanColCons_lemma]);
val ⊆_cleanColCons_lemma1 = save_pop_thm"⊆_cleanColCons_lemma1";

```

HOL output

```

⊆_cleanColCons_lemma1 =
| ⊢ ∀ c1 c2 t col
|   • c1 dominates c2
|     ⇒ col ∈ Snd (cleanColCons c2 t)
|     ⇒ col ∈ Snd (cleanColCons c1 t)

```

The main result about *CleanRow*.

SML

```

push_goal([],Γ∀ c1 c2 • c1 dominates c2
  ⇒ ∀ t1 t2 r1 r2 •
    (cleanColCons c1 t1 = cleanColCons c1 t2
      ∧ cleanColCons c2 t1 = cleanColCons c2 t2
      ∧ cleanRow c1 (Snd (cleanColCons c1 t1)) r1
        = cleanRow c1 (Snd (cleanColCons c1 t2)) r2)
    ⇒
      cleanRow c2 (Snd (cleanColCons c2 t1)) r1
        = cleanRow c2 (Snd (cleanColCons c2 t2)) r2∇);
a(REPEAT ∀_tac THEN strip_tac);
a(strip_asm_tac (list_∇_elim[Γc1∇,Γc2∇]replaceData_lemma));
a(rewrite_tac[cleanRow_def,row_components,get_specΓMkRow∇,filterRow_def,
  rel_ext_clauses,∠_thm,r_∅_r_thm,graph_thm] THEN REPEAT strip_tac);
(* *** Goal "1" *** *)
a(strip_asm_tac (list_∇_elim[Γc1∇,Γc2∇,Γt1∇,Γc∇]⊆_cleanColCons_lemma1));
a(DROP_NTH_ASM_T 6 (asm_tac o list_∇_elim[Γx∇,ΓreplaceData c1 z∇]));
a(LEMMA_TΓ∃ z'
  • ((∃ c • c ∈ Snd (cleanColCons c1 t1) ∧ CS_posn c = x)
    ∧ (x, z') ∈ R_data r1)
    ∧ replaceData c1 z = replaceData c1 z'∇asm_tac);
(* *** Goal "1.1" *** *)
a(∃_tacΓz∇ THEN asm_prove_tac[]);

```

SML

```

(* *** Goal "1.2" *** *)
a(DROP_NTH_ASM_T 2 ante_tac THEN POP_ASM_T rewrite_thm_tac);
a  $\Rightarrow$  _tac;
a(DROP_NTH_ASM_T 13(strip_asm_tac o list_ $\forall$ _elim[ $\Gamma z^\neg$ , $\Gamma z'^\neg$ ]));
a( $\exists$ _tac $\Gamma z'^\neg$  THEN asm_rewrite_tac[]);
a( $\exists$ _tac $\Gamma c^\neg$  THEN asm_rewrite_tac[]);
a(DROP_NTH_ASM_T 10 ante_tac THEN asm_rewrite_tac[]);

```

SML

```

(* *** Goal "2" *** *)
a(strip_asm_tac (list_ $\forall$ _elim[ $\Gamma c_1^\neg$ , $\Gamma c_2^\neg$ , $\Gamma t_2^\neg$ , $\Gamma c^\neg$ ] $\sqsubseteq$ _cleanColCons_lemma1));
a(DROP_NTH_ASM_T 6 (asm_tac o list_ $\forall$ _elim[ $\Gamma x^\neg$ , $\Gamma replaceData c_1 z^\neg$ ]));
a(LEMMA_T $\Gamma \exists z'$ 
  • (( $\exists c \bullet c \in Snd$  (cleanColCons  $c_1 t_2$ )  $\wedge CS\_posn c = x$ )
     $\wedge (x, z') \in R\_data r_2$ )
     $\wedge replaceData c_1 z = replaceData c_1 z'^\neg$ asm_tac);
(* *** Goal "2.1" *** *)
a( $\exists$ _tac $\Gamma z^\neg$  THEN asm_prove_tac[]);

```

SML

```

(* *** Goal "2.2" *** *)
a(DROP_NTH_ASM_T 2 ante_tac THEN POP_ASM_T rewrite_thm_tac);
a  $\Rightarrow$  _tac;
a(DROP_NTH_ASM_T 13(strip_asm_tac o list_ $\forall$ _elim[ $\Gamma z^\neg$ , $\Gamma z'^\neg$ ]));
a( $\exists$ _tac $\Gamma z'^\neg$  THEN asm_rewrite_tac[]);
a( $\exists$ _tac $\Gamma c^\neg$  THEN asm_rewrite_tac[]);
val cleanRow_lemma = save_pop_thm"cleanRow_lemma";

```

HOL output

```

cleanRow_lemma =
|  $\vdash \forall c_1 c_2$ 
  •  $c_1$  dominates  $c_2$ 
     $\Rightarrow (\forall t_1 t_2 r_1 r_2$ 
      • cleanColCons  $c_1 t_1 = cleanColCons c_1 t_2$ 
         $\wedge cleanColCons c_2 t_1 = cleanColCons c_2 t_2$ 
         $\wedge cleanRow c_1 (Snd (cleanColCons c_1 t_1)) r_1$ 
          = cleanRow  $c_1 (Snd (cleanColCons c_1 t_2)) r_2$ 
         $\Rightarrow cleanRow c_2 (Snd (cleanColCons c_2 t_1)) r_1$ 
          = cleanRow  $c_2 (Snd (cleanColCons c_2 t_2)) r_2$ )

```

**4.2.4 cleanTable Lemma**

SML

```

| push_goal([],Γ∀ c1 c2 • c1 dominates c2 ⇒ ∀ t1 t2 •
|   cleanTable c1 t1 = cleanTable c1 t2 ⇒ cleanTable c2 t1 = cleanTable c2 t2∇);
| a(rewrite_tac[cleanTable_def] THEN REPEAT ∀_tac THEN strip_tac THEN REPEAT ∀_tac);
| a(cases_tacΓc2 dominates TS_class t1∇ THEN cases_tac Γc2 dominates TS_class t2∇
|   THEN asm_rewrite_tac[]);

```

SML

```

| (* *** Goal "1" *** *)
| a(lemma_tacΓc1 dominates TS_class t1∇ THEN_LIST
|   [contr_tac THEN basic_res_tac4 2 [1][3,4][dominates_trans],asm_rewrite_tac[]]);
| a(lemma_tacΓc1 dominates TS_class t2∇ THEN_LIST
|   [contr_tac THEN basic_res_tac4 2 [1][3,5][dominates_trans],asm_rewrite_tac[]]);
| a(rewrite_tac[tab_components,get_specΓMkTableSpec∇]);
| a(strip_tac THEN asm_rewrite_tac[]);
| a(strip_asm_tac (list_∇_elimΓcleanColCons c1 t1∇,ΓcleanColCons c1 t2∇pair_eq_thm1));
| a(rewrite_tac[taut_ruleΓ∀ a b c •(a ∧ (b ∧ c)) = ((b ∧ a) ∧ c)∇,pair_eq_thm1]);
| a(DROP_NTH_ASM_T 4 (fn _ => id_tac)
|   THEN DROP_NTH_ASM_T 3 (fn _ => id_tac));
| a(strip_asm_tac(list_∇_elimΓc1∇,Γc2∇,Γt1∇,Γt2∇cleanColCons_lemma));
| a(∧_tac THEN_LIST[asm_rewrite_tac[],id_tac]);

```

SML

```

| a(DROP_NTH_ASM_T 3 ante_tac);
| a(lemma_tacΓ∃ l • TS_rows t1 = l∇THEN_LIST
|   [prove_∃_tac,POP_ASM_T rewrite_thm_tac]);
| a(lemma_tacΓ∃ l1 • TS_rows t2 = l1∇THEN_LIST
|   [prove_∃_tac,POP_ASM_T rewrite_thm_tac]);
| a(intro_∇_tac(Γl1∇,Γl':Row LIST∇));
| a(rewrite_tac[cleanRows_def]);
| a(LIST_INDUCTION_TΓl∇asm_tac);
| (* *** Goal "1.1" *** *)
| a(rewrite_tac[map_def,map_null_thm]);
| a ∀_tac;
| a(LIST_INDUCTION_TΓl'∇asm_tac);

```



SML

```

(* *** Goal "1.1.1" *** *)
a(rewrite_tac[|-def]);
(* *** Goal "1.1.2" *** *)
a(rewrite_tac[|-def]);
a strip_tac;
a(cases_tacΓc2 dominates R_exist x∇ THEN asm_rewrite_tac[]);
(* *** Goal "1.1.2.1" *** *)
a(lemma_tacΓc1 dominates R_exist x∇ THEN_LIST
  [contr_tac THEN basic_res_tac4 2 [1][2,11]] [dominates_trans],asm_rewrite_tac[]);
(* *** Goal "1.1.2.2" *** *)
a(cases_tacΓc1 dominates R_exist x∇ THEN asm_rewrite_tac[]);

```

SML

```

(* *** Goal "1.2" *** *)
a(rewrite_tac[∀_reorder_convΓ∀ l' x
  • Map
    (cleanRow c1 (Snd (cleanColCons c1 t1)))
    (Cons x l | {r|c1 dominates R_exist r})
  = Map
    (cleanRow c1 (Snd (cleanColCons c1 t2)))
    (l' | {r|c1 dominates R_exist r})
  ⇒ Map
    (cleanRow c2 (Snd (cleanColCons c2 t1)))
    (Cons x l | {r|c2 dominates R_exist r})
  = Map
    (cleanRow c2 (Snd (cleanColCons c2 t2)))
    (l' | {r|c2 dominates R_exist r})∇∀ x l'
  • Map
    (cleanRow c1 (Snd (cleanColCons c1 t1)))
    (Cons x l | {r|c1 dominates R_exist r})
  = Map
    (cleanRow c1 (Snd (cleanColCons c1 t2)))
    (l' | {r|c1 dominates R_exist r})
  ⇒ Map
    (cleanRow c2 (Snd (cleanColCons c2 t1)))
    (Cons x l | {r|c2 dominates R_exist r})
  = Map
    (cleanRow c2 (Snd (cleanColCons c2 t2)))
    (l' | {r|c2 dominates R_exist r})∇]);

```

SML

```

a  $\forall$ _tac;
a(LIST_INDUCTION_TΓ l'∇ asm_tac);
(* *** Goal "1.2.1" *** *)
a(rewrite_tac[ $\lambda$ _def] THEN  $\forall$ _tac);
a(cases_tacΓ c2 dominates R_exist x∇ THEN TOP_ASM_T rewrite_thm_tac);
(* *** Goal "1.2.1.1" *** *)
a(lemma_tacΓ c1 dominates R_exist x∇ THEN_LIST
  [contr_tac THEN basic_res_tac4 2 [1][2,4]] [[dominates_trans],
  TOP_ASM_T rewrite_thm_tac]);
a(rewrite_tac[map_def]);
(* *** Goal "1.2.1.2" *** *)
a(cases_tacΓ c1 dominates R_exist x∇ THEN TOP_ASM_T rewrite_thm_tac);
(* *** Goal "1.2.1.2.1" *** *)
a(rewrite_tac[map_def]);
(* *** Goal "1.2.1.2.2" *** *)
a(DROP_NTH_ASM_T 3(ante_tac o  $\forall$ _elimΓ []:Row LIST∇));
a(rewrite_tac[map_def]);

```

SML

```

(* *** Goal "1.2.2" *** *)
a(rewrite_tac[ $\lambda$ _def]);
a(REPEAT  $\forall$ _tac);
a(cases_tacΓ c2 dominates R_exist x∇
  THEN TOP_ASM_T rewrite_thm_tac THEN cases_tacΓ c2 dominates R_exist x'∇
  THEN TOP_ASM_T rewrite_thm_tac);
(* *** Goal "1.2.2.1" *** *)
a(lemma_tacΓ c1 dominates R_exist x∇ THEN_LIST
  [contr_tac THEN basic_res_tac4 2 [1][3,14]] [[dominates_trans],
  TOP_ASM_T rewrite_thm_tac]);
a(lemma_tacΓ c1 dominates R_exist x'∇ THEN_LIST
  [contr_tac THEN basic_res_tac4 2 [1][3,14]] [[dominates_trans],
  TOP_ASM_T rewrite_thm_tac]);
a(rewrite_tac[map_def]);

```

SML

```

a(DROP_NTH_ASM_T 6 (asm_tac o  $\forall$ _elim $\ulcorner$ l' $\urcorner$ ));
a(strip_asm_tac(list_ $\forall$ _elim $\ulcorner$ c1 $\urcorner$ , $\ulcorner$ c2 $\urcorner$ ]cleanRow_lemma));
a(POP_ASM_T (ante_tac o list_ $\forall$ _elim $\ulcorner$ t1 $\urcorner$ , $\ulcorner$ t2 $\urcorner$ , $\ulcorner$ x' $\urcorner$ , $\ulcorner$ x $\urcorner$ ));
a(asm_rewrite_tac[] THEN  $\Rightarrow$ _T asm_tac);
a strip_tac;
a(DROP_NTH_ASM_T 4 ante_tac THEN DROP_NTH_ASM_T 3 ante_tac
  THEN asm_rewrite_tac[] THEN REPEAT strip_tac);
(* *** Goal "1.2.2.2" *** *)
a(lemma_tac $\ulcorner$ c1 dominates R_exist x $\urcorner$  THEN_LIST
  [contr_tac THEN basic_res_tac4 2 [1][3,14]][[dominates_trans],
  TOP_ASM_T rewrite_thm_tac]);
a(cases_tac  $\ulcorner$ c1 dominates R_exist x' $\urcorner$  THEN asm_rewrite_tac[]);
(* *** Goal "1.2.2.2.1" *** *)
a(rewrite_tac[map_def] THEN strip_tac);
a(DROP_NTH_ASM_T 2
  (ante_tac o rewrite_rule[cleanRow_def,row_components,get_spec $\ulcorner$ MkRow $\urcorner$ ]));
a strip_tac;
a(DROP_NTH_ASM_T 6 ante_tac THEN asm_rewrite_tac[]);

```

SML

```

(* *** Goal "1.2.2.2.2" *** *)
a(rewrite_tac[map_def] THEN strip_tac);
a(DROP_NTH_ASM_T 7(ante_tac o  $\forall$ _elim $\ulcorner$ Cons x l' $\urcorner$ ));
a(asm_rewrite_tac[map_def,|-def]);
(* *** Goal "1.2.2.3" *** *)
a(lemma_tac $\ulcorner$ c1 dominates R_exist x' $\urcorner$  THEN_LIST
  [contr_tac THEN basic_res_tac4 2 [1][2,14]][[dominates_trans],
  TOP_ASM_T rewrite_thm_tac]);
a(cases_tac  $\ulcorner$ c1 dominates R_exist x' $\urcorner$  THEN asm_rewrite_tac[]);
(* *** Goal "1.2.2.3.1" *** *)
a(rewrite_tac[map_def] THEN strip_tac);
a(DROP_NTH_ASM_T 2(ante_tac o
  rewrite_rule[cleanRow_def,row_components,get_spec $\ulcorner$ MkRow $\urcorner$ ]));
a strip_tac;
a(DROP_NTH_ASM_T 6 ante_tac THEN asm_rewrite_tac[]);

```

SML

```

(* *** Goal "1.2.2.3.2" *** *)
a(rewrite_tac[map_def] THEN strip_tac);
a(DROP_NTH_ASM_T 6(ante_tac o  $\forall$ _elim $\lceil$ x' $\rceil$ ));
a(asm_rewrite_tac[map_def,  $\lceil$ -def]);
(* *** Goal "1.2.2.4" *** *)
a(cases_tac  $\lceil$ c1 dominates R_exist x' $\rceil$  THEN cases_tac  $\lceil$ c1 dominates R_exist x' $\rceil$ 
  THEN asm_rewrite_tac[]);
(* *** Goal "1.2.2.4.1" *** *)
a(rewrite_tac[map_def] THEN strip_tac);
a(DROP_NTH_ASM_T 8(ante_tac o  $\forall$ _elim $\lceil$ l' $\rceil$ ) THEN asm_rewrite_tac[]);

```

SML

```

(* *** Goal "1.2.2.4.2" *** *)
a(rewrite_tac[map_def] THEN strip_tac);
a(DROP_NTH_ASM_T 7(ante_tac o  $\forall$ _elim $\lceil$ Cons x l' $\rceil$ ));
a(asm_rewrite_tac[map_def,  $\lceil$ -def]);
(* *** Goal "1.2.2.4.3" *** *)
a(rewrite_tac[map_def] THEN strip_tac);
a(DROP_NTH_ASM_T 6(ante_tac o  $\forall$ _elim $\lceil$ x' $\rceil$ ));
a(asm_rewrite_tac[map_def,  $\lceil$ -def]);
(* *** Goal "1.2.2.4.4" *** *)
a(strip_tac THEN DROP_NTH_ASM_T 7(ante_tac o  $\forall$ _elim $\lceil$ l' $\rceil$ ) THEN asm_rewrite_tac[]);

```

SML

```

(* *** Goal "2" *** *)
a(lemma_tac $\lceil$ c1 dominates TS_class t1 $\rceil$  THEN_LIST
  [contr_tac THEN basic_res_tac4 2 [1][3,4][][][dominates_trans],asm_rewrite_tac[]]);
a(cases_tac $\lceil$ c1 dominates TS_class t2 $\rceil$  THEN asm_rewrite_tac[]);
(* *** Goal "2.1" *** *)
a(rewrite_tac[tab_components,get_spec $\lceil$ MkTableSpec $\rceil$ ]);
a strip_tac;
a(DROP_NTH_ASM_T 9 ante_tac THEN asm_rewrite_tac[]);
(* *** Goal "2.2" *** *)
a(rewrite_tac[tab_components,get_spec $\lceil$ MkTableSpec $\rceil$ ]);
a(strip_tac THEN asm_rewrite_tac[]);
a(rewrite_tac[taut_rule $\lceil$  $\forall$  a b c • (a  $\wedge$  (b  $\wedge$  c)) = ((b  $\wedge$  a)  $\wedge$  c) $\rceil$ ]);
a strip_tac;

```

SML

```

(* *** Goal "2.2.1" *** *)
a(ante_tac(list_∇_elim[Γ c1 ⊃, Γ c2 ⊃, Γ t1 ⊃] ⊆_cleanColCons_lemma));
a(asm_rewrite_tac[⊆_∅_thm] THEN REPEAT strip_tac);
(* *** Goal "2.2.2" *** *)
a(POP_ASM_T ante_tac THEN
  rewrite_tac[cleanRows_def, map_null_thm, ⊃_thm4a, sets_ext_clauses]);
a(REPEAT strip_tac);

```

SML

```

(* *** Goal "2.2.2.1" *** *)
a(lemma_tac[Γ c1 dominates R_exist x ⊃ THEN_LIST
  [contr_tac THEN basic_res_tac4 2 [1][3,13]] [[dominates_trans], id_tac]);
a(DROP_NTH_ASM_T 4 (strip_asm_tac o ∇_elim[Γ x ⊃]));
(* *** Goal "2.2.2.2" *** *)
a(lemma_tac[Γ c1 dominates R_exist x' ⊃ THEN_LIST
  [contr_tac THEN basic_res_tac4 2 [1][2,13]] [[dominates_trans], basic_res_tac1 1 []]);

```

SML

```

(* *** Goal "3" *** *)
a(lemma_tac[Γ c1 dominates TS_class t2 ⊃ THEN_LIST
  [contr_tac THEN basic_res_tac4 2 [1][2,4]] [[dominates_trans], asm_rewrite_tac []]);
a(cases_tac[Γ c1 dominates TS_class t1 ⊃ THEN asm_rewrite_tac []]);
(* *** Goal "3.1" *** *)
a(rewrite_tac[tab_components, get_spec[Γ MkTableSpec ⊃]);
a strip_tac;
a(DROP_NTH_ASM_T 9 ante_tac THEN asm_rewrite_tac []);

```

SML

```

(* *** Goal "3.2" *** *)
a(rewrite_tac[tab_components, get_spec[Γ MkTableSpec ⊃]);
a strip_tac;
a(DROP_NTH_ASM_T 5(rewrite_thm_tac o eq_sym_rule));
a(DROP_NTH_ASM_T 4(rewrite_thm_tac o eq_sym_rule));
a(rewrite_tac[taut_rule[∇ a b c • (a ∧ (b ∧ c)) = ((b ∧ a) ∧ c) ⊃] THEN strip_tac);
(* *** Goal "3.2.1" *** *)
a(ante_tac(list_∇_elim[Γ c1 ⊃, Γ c2 ⊃, Γ t2 ⊃] ⊆_cleanColCons_lemma));
a(DROP_NTH_ASM_T 3 (rewrite_thm_tac o eq_sym_rule));
a(DROP_NTH_ASM_T 2 (rewrite_thm_tac o eq_sym_rule));
a(asm_rewrite_tac[⊆_∅_thm1] THEN REPEAT strip_tac);

```

SML

```

(* *** Goal "3.2.2" *** *)
a(POP_ASM_T ante_tac THEN
  rewrite_tac[cleanRows_def, map_null_thm, |-_thm4a, sets_ext_clauses]);
a(REPEAT strip_tac);
(* *** Goal "3.2.2.1" *** *)
a(lemma_tacΓ c1 dominates R_exist x∇ THEN_LIST
  [contr_tac THEN basic_res_tac4 2 [1][3,11][[dominates_trans], id_tac]);
a(DROP_NTH_ASM_T 4 (strip_asm_tac o ∇_elimΓ x∇));

```

SML

```

(* *** Goal "3.2.2.2" *** *)
a(lemma_tacΓ c1 dominates R_exist x'∇ THEN_LIST
  [contr_tac THEN basic_res_tac4 2 [1][2,11][[dominates_trans], basic_res_tac1 1 []]);
(* *** Goal "4" *** *)
a(cases_tacΓ c1 dominates TS_class t1∇ THEN cases_tacΓ c1 dominates TS_class t2∇
  THEN asm_rewrite_tac[tab_components, get_specΓ MkTableSpec∇]
  THEN strip_tac THEN asm_rewrite_tac[]);
val cleanTable_lemma = save_pop_thm "cleanTable_lemma";

```

HOL output

```

cleanTable_lemma =
| ⊢ ∇ c1 c2
| • c1 dominates c2
|   ⇒ (∇ t1 t2
|     • cleanTable c1 t1 = cleanTable c1 t2
|       ⇒ cleanTable c2 t1 = cleanTable c2 t2)

```

#### 4.2.5 cleanDirectory Lemma

SML

```

push_goal([],Γ ∇ c1 c2 • c1 dominates c2 ⇒ ∇ d1 d2 •
  cleanDirectory c1 d1 = cleanDirectory c1 d2
  ⇒ cleanDirectory c2 d1 = cleanDirectory c2 d2∇);
a(REPEAT strip_tac);
a(POP_ASM_T ante_tac THEN rewrite_tac[cleanDirectory_def]);
a(cases_tacΓ c2 dominates Dir_class d1∇ THEN cases_tacΓ c2 dominates Dir_class d2∇
  THEN asm_rewrite_tac[]);

```

SML

```

(* *** Goal "1" *** *)
a(lemma_tacΓc1 dominates Dir_class d1∇ THEN_LIST
  [contr_tac THEN basic_res_tac4 2 [1][3,4][][dominates_trans],asm_rewrite_tac[]]);
a(lemma_tacΓc1 dominates Dir_class d2∇ THEN_LIST
  [contr_tac THEN basic_res_tac4 2 [1][3,5][][dominates_trans],asm_rewrite_tac[]]);
a(rewrite_tac[dir_components,get_specΓMkDirectory∇] THEN REPEAT strip_tac);
a(strip_asm_tac(list_∇_elimΓc1∇,Γc2∇]cleanTable_lemma));
a(ante_tac (list_∇_elimΓcleanTable c1∇,ΓcleanTable c2∇,
  ΓDir_tables d1∇,ΓDir_tables d2∇]rel_thm) THEN asm_rewrite_tac[]);

```

SML

```

(* *** Goal "2" *** *)
a(lemma_tacΓc1 dominates Dir_class d1∇ THEN_LIST
  [contr_tac THEN basic_res_tac4 2 [1][3,4][][dominates_trans],asm_rewrite_tac[]]);
a(cases_tacΓc1 dominates Dir_class d2∇ THEN asm_rewrite_tac[]);
(* *** Goal "2.1" *** *)
a(rewrite_tac[dir_components,get_specΓMkDirectory∇] THEN REPEAT strip_tac);
a(DROP_NTH_ASM_T 7 ante_tac THEN asm_rewrite_tac[]);
(* *** Goal "2.2" *** *)
a(rewrite_tac[dir_components,get_specΓMkDirectory∇] THEN REPEAT strip_tac);
a(DROP_NTH_ASM_T 3 ante_tac THEN rewrite_tac[graph_thm,r_∅_r_thm,rel_ext_clauses]);
a(REPEAT strip_tac);
a(DROP_NTH_ASM_T 2 (strip_asm_tac o list_∇_elimΓx∇,ΓcleanTable c1 z∇));
a(POP_ASM_T (strip_asm_tac o ∇_elimΓz∇));

```

SML

```

(* *** Goal "3" *** *)
a(lemma_tacΓc1 dominates Dir_class d2∇ THEN_LIST
  [contr_tac THEN basic_res_tac4 2 [1][2,4][][dominates_trans],asm_rewrite_tac[]]);
a(cases_tacΓc1 dominates Dir_class d1∇ THEN asm_rewrite_tac[]);
(* *** Goal "3.1" *** *)
a(rewrite_tac[dir_components,get_specΓMkDirectory∇] THEN REPEAT strip_tac);
a(DROP_NTH_ASM_T 7 ante_tac THEN asm_rewrite_tac[]);
(* *** Goal "3.2" *** *)
a(rewrite_tac[dir_components,get_specΓMkDirectory∇] THEN REPEAT strip_tac);
a(DROP_NTH_ASM_T 3 ante_tac THEN rewrite_tac[graph_thm,r_∅_r_thm,rel_ext_clauses]);

```

SML

```

a(REPEAT strip_tac);
a(DROP_NTH_ASM_T 2 (strip_asm_tac o list_∇_elim[⌈x⌋,⌈cleanTable c1 z⌋]));
a(POP_ASM_T (strip_asm_tac o ∇_elim⌈z⌋));
(* *** Goal "4" *** *)
a(cases_tac⌈c1 dominates Dir_class d1⌋ THEN cases_tac⌈c1 dominates Dir_class d2⌋
  THEN asm_rewrite_tac[dir_components,get_spec⌈MkDirectory⌋]
  THEN REPEAT strip_tac);
val cleanDirectory_lemma = save_pop_thm"cleanDirectory_lemma";

```

HOL output

```

cleanDirectory_lemma =
| ⊢ ∇ c1 c2
| • c1 dominates c2
|   ⇒ (∇ d1 d2
|     • cleanDirectory c1 d1 = cleanDirectory c1 d2
|     ⇒ cleanDirectory c2 d1 = cleanDirectory c2 d2)

```

#### 4.2.6 hideR Lemma

First, two lemmas about directories.

SML

```

push_goal([],⌈∇ c1 c2 s x y • c1 dominates c2 ⇒
  ((x,y) ∈ s ▷ {dir|c2 dominates Dir_exist dir}
  ⇒ (x,y) ∈ s ▷ {dir|c1 dominates Dir_exist dir})⌋);
a(rewrite_tac[▷_thm] THEN REPEAT strip_tac);
a(contr_tac THEN basic_res_tac4 2 [1][3,4][[]][dominates_trans]);
val ⊆_dir_lemma = save_pop_thm"⊆_dir_lemma";

```

HOL output

```

⊆_dir_lemma =
| ⊢ ∇ c1 c2 s x y
| • c1 dominates c2
|   ⇒ (x, y) ∈ s ▷ {dir|c2 dominates Dir_exist dir}
|   ⇒ (x, y) ∈ s ▷ {dir|c1 dominates Dir_exist dir}

```



SML

```

push_goal([],Γ∀ c c1 s s1 x z z1
  • (cleanDirectory c1 z = cleanDirectory c1 z1
    ∧ (x, z) ∈ s ▷ {dir|c1 dominates Dir_exist dir}
    ∧ (x, z1) ∈ s1 ▷ {dir|c dominates Dir_exist dir})
    ⇒ (x, z1) ∈ s1 ▷ {dir|c1 dominates Dir_exist dir}¬);
a(rewrite_tac[cleanDirectory_def,dir_components,get_specΓMkDirectory¬,▷_thm]
  THEN REPEAT strip_tac);
a(DROP_NTH_ASM_T 4 ante_tac THEN asm_rewrite_tac[]);
val dir_class_preserved_lemma = save_pop_thm"dir_class_preserved_lemma";

```

HOL output

```

dir_class_preserved_lemma =
| ⊢ ∀ c c1 s s1 x z z1
| • cleanDirectory c1 z = cleanDirectory c1 z1
|   ∧ (x, z) ∈ s ▷ {dir|c1 dominates Dir_exist dir}
|   ∧ (x, z1) ∈ s1 ▷ {dir|c dominates Dir_exist dir}
|   ⇒ (x, z1) ∈ s1 ▷ {dir|c1 dominates Dir_exist dir}

```

Now the main result of this section that *hideR* is a member of *secureHideR*.

SML

```

push_goal([],ΓhideR ∈ secureHideR¬);
a(rewrite_tac[get_specΓhideR¬,get_specΓsecureHideR¬]);
a(REPEAT strip_tac);
a(strip_asm_tac(list_∀_elim[Γc1¬,Γc2¬]cleanDirectory_lemma));
a(DROP_NTH_ASM_T 3 ante_tac);
a(rewrite_tac[r_9_r_thm,rel_ext_clauses,graph_thm]THEN REPEAT strip_tac);
(* *** Goal "1" *** *)
a(strip_asm_tac(list_∀_elim[Γc1¬,Γc2¬,Γs1¬,Γx¬,Γz¬]⊆_dir_lemma));
a(DROP_NTH_ASM_T 4 (asm_tac o list_∀_elim[Γx¬,ΓcleanDirectory c1 z¬]));
a(LEMMA_TΓ∃ z'
  • (x, z') ∈ s1 ▷ {dir|c1 dominates Dir_exist dir}
    ∧ cleanDirectory c1 z = cleanDirectory c1 z'¬asm_tac);
(* *** Goal "1.1" *** *)
a(∃_tacΓz¬ THEN asm_prove_tac[]);

```

SML

```

(* *** Goal "1.2" *** *)
a(DROP_NTH_ASM_T 2 ante_tac THEN POP_ASM_T rewrite_thm_tac);
a  $\Rightarrow$  _tac;
a(DROP_NTH_ASM_T 6(strip_asm_tac o list_ $\forall$ _elim[ $\Gamma$  z $\neg$ , $\Gamma$  z' $\neg$ ]));
a( $\exists$ _tac $\Gamma$  z' $\neg$  THEN asm_rewrite_tac[]);
a(strip_asm_tac (list_ $\forall$ _elim[ $\Gamma$  c1 $\neg$ , $\Gamma$  c2 $\neg$ , $\Gamma$  s1 $\neg$ , $\Gamma$  s2 $\neg$ , $\Gamma$  x $\neg$ , $\Gamma$  z $\neg$ , $\Gamma$  z' $\neg$ ]
  dir_class_preserved_lemma));
(* *** Goal "2" *** *)
a(strip_asm_tac (list_ $\forall$ _elim[ $\Gamma$  c1 $\neg$ , $\Gamma$  c2 $\neg$ , $\Gamma$  s2 $\neg$ , $\Gamma$  x $\neg$ , $\Gamma$  z $\neg$ ] $\subseteq$ _dir_lemma));
a(DROP_NTH_ASM_T 4 (asm_tac o list_ $\forall$ _elim[ $\Gamma$  x $\neg$ , $\Gamma$  cleanDirectory c1 z $\neg$ ]));
a(LEMMA_T $\Gamma$  $\exists$  z'
  • (x, z')  $\in$  s2  $\triangleright$  {dir | c1 dominates Dir_exist dir}
     $\wedge$  cleanDirectory c1 z = cleanDirectory c1 z' $\neg$ asm_tac);
(* *** Goal "2.1" *** *)
a( $\exists$ _tac $\Gamma$  z $\neg$  THEN asm_prove_tac[]);

```

SML

```

(* *** Goal "2.2" *** *)
a(DROP_NTH_ASM_T 2 ante_tac THEN POP_ASM_T rewrite_thm_tac);
a  $\Rightarrow$  _tac;
a(DROP_NTH_ASM_T 6(strip_asm_tac o list_ $\forall$ _elim[ $\Gamma$  z $\neg$ , $\Gamma$  z' $\neg$ ]));
a( $\exists$ _tac $\Gamma$  z' $\neg$  THEN asm_rewrite_tac[]);
a(strip_asm_tac (list_ $\forall$ _elim[ $\Gamma$  c1 $\neg$ , $\Gamma$  c2 $\neg$ , $\Gamma$  s2 $\neg$ , $\Gamma$  s1 $\neg$ , $\Gamma$  x $\neg$ , $\Gamma$  z $\neg$ , $\Gamma$  z' $\neg$ ]
  dir_class_preserved_lemma));
val secureHideR_lemma = save_pop_thm"secureHideR_lemma";

```

HOL output

```

secureHideR_lemma =  $\vdash$  hideR  $\in$  secureHideR

```

### 4.3 Proof that the Invariant on the State is Maintained

We need a series of auxiliary results of the form *replaceData*, etc., is partial.

SML

```

push_goal([], $\Gamma$  $\forall$  clear • Graph (cleanDirectory clear)  $\in$  Functional $\neg$ );
a(rewrite_tac[cleanDirectory_def,graph_thm,functional_def,dir_components,
  get_spec $\Gamma$  MkDirectory $\neg$ ] THEN REPEAT strip_tac THEN asm_rewrite_tac[]);
val cleanDir_fun_thm = save_pop_thm"cleanDir_fun_thm";

```

HOL output

```

cleanDir_fun_thm =  $\vdash$   $\forall$  clear • Graph (cleanDirectory clear)  $\in$  Functional

```

SML

```

| push_goal([],⌈∀ clear • Graph (replaceData clear) ∈ Functional⌋);
| a(rewrite_tac[replaceData_def,graph_thm,functional_def,data_components,get_spec⌈MkData⌋]
|   THEN REPEAT strip_tac THEN asm_rewrite_tac[]);
| val replaceData_fun_thm = save_pop_thm"replaceData_fun_thm";

```

HOL output

```

| replaceData_fun_thm = ⊢ ∀ clear • Graph (replaceData clear) ∈ Functional

```

SML

```

| push_goal([],⌈∀ clear tab • TS_cons tab ∈ Functional
|   ⇒ Fst (cleanColCons clear tab) ∈ Functional⌋);
| a(rewrite_tac[cleanColCons_def,functional_def,>-thm] THEN REPEAT strip_tac);
| a(DROP_NTH_ASM_T 5 (strip_asm_tac o list_∀_elim[⌈x⌋,⌈w⌋,⌈z⌋]));
| val cleanColCons_fun_thm = save_pop_thm"cleanColCons_fun_thm";

```

HOL output

```

| cleanColCons_fun_thm =
| ⊢ ∀ clear tab
|   • TS_cons tab ∈ Functional ⇒ Fst (cleanColCons clear tab) ∈ Functional

```

SML

```

| push_goal([],⌈∀ clear dir • Dir_tables dir ∈ Functional
|   ⇒ Dir_tables (cleanDirectory clear dir) ∈ Functional⌋);
| a(rewrite_tac[cleanDirectory_def,functional_def] THEN REPEAT strip_tac);
| a(POP_ASM_T ante_tac THEN POP_ASM_T ante_tac
|   THEN rewrite_tac[get_spec⌈MkDirectory⌋]);
| a(cases_tac⌈clear dominates Dir_class dir⌋ THEN asm_rewrite_tac[]);
| a(rewrite_tac[r_g-r_thm,graph_thm]);
| a(REPEAT strip_tac);
| a(DROP_NTH_ASM_T 6 (strip_asm_tac o list_∀_elim[⌈x⌋,⌈z'⌋,⌈z''⌋]));
| a(asm_rewrite_tac[]);
| val tables_fun_thm = save_pop_thm"tables_fun_thm";

```

HOL output

```

| tables_fun_thm =
| ⊢ ∀ clear dir
|   • Dir_tables dir ∈ Functional
|     ⇒ Dir_tables (cleanDirectory clear dir) ∈ Functional

```

Now the main result of this section that *hideR* preserves the invariant on states.

SML

```

| push_goal([],⊢∀ clear s • isState s ⇒ isState(hideR(clear,s))⊣);
| a(rewrite_tac[get_spec⊢hideR⊣] THEN REPEAT ∀_tac);
| a(rewrite_tac[get_spec⊢isState⊣,get_spec⊢StateS⊣,↔_def,∩_def] THEN strip_tac);
| a(lemma_tac⊢(s ▷ {dir|clear dominates Dir_exist dir}) % Graph (cleanDirectory clear)
|   ∈ Functional⊣);
| (* *** Goal "1" *** *)
| a(strip_asm_tac(list_∀_elim⊢{dir|clear dominates Dir_exist dir}⊣,⊢s⊣)▷_fun_thm));
| a(strip_asm_tac(∀_elim⊢clear⊣cleanDir_fun_thm));
| a(strip_asm_tac(list_∀_elim⊢(s ▷ {dir|clear dominates Dir_exist dir})⊣,
|   ⊢cleanDirectory clear⊣)%_fun_thm));

```

SML

```

| (* *** Goal "2" *** *)
| a(TOP_ASM_T rewrite_thm_tac);
| a(DROP_NTH_ASM_T 3 ante_tac THEN rewrite_tac[↔_def,get_spec⊢IdeL⊣,
|   get_spec⊢DirectoryS⊣,∩_def,×_def,get_spec⊢Universe⊣,dom_def,
|   ▷_thm,r_%_r_thm,graph_thm,rel_ext_clauses,get_spec⊢$P⊣]);
| a(REPEAT strip_tac);
| a(asm_rewrite_tac[]);
| a(DROP_NTH_ASM_T 4(strip_asm_tac o list_∀_elim⊢x⊣,⊢z⊣));
| a(POP_ASM_T ante_tac THEN rewrite_tac[↔_def,∩_def] THEN strip_tac);
| a(strip_asm_tac(list_∀_elim⊢clear⊣,⊢z⊣)tables_fun_thm));
| a(POP_ASM_T rewrite_thm_tac);
| a(DROP_NTH_ASM_T 2 ante_tac);

```

SML

```

| a(DROP_NTH_ASM_T 2 ante_tac
|   THEN rewrite_tac[cleanDirectory_def,get_spec⊢MkDirectory⊣]
|   THEN cases_tac⊢clear dominates Dir_class z⊣
|   THEN asm_rewrite_tac[↔_def,×_def] THEN strip_tac);
| a(rewrite_tac[∈_in_clauses,rel_ext_clauses,get_spec⊢TableSpecS⊣,
|   get_spec⊢Ide⊣,↔_def,∩_def,r_%_r_thm,graph_thm,cleanTable_def]
|   THEN conv_tac(MAP_C let_conv) THEN strip_tac THEN strip_tac
|   THEN strip_tac THEN strip_tac);
| a(DROP_NTH_ASM_T 3(strip_asm_tac o list_∀_elim⊢x'⊣,⊢z'⊣));
| a(DROP_NTH_ASM_T 4 ante_tac THEN cases_tac⊢clear dominates TS_class z'⊣
|   THEN asm_rewrite_tac[tab_components,get_spec⊢MkTableSpec⊣]
|   THEN strip_tac
|   THEN asm_rewrite_tac[↔_def,×_def,prove_rule[functional_def]⊢{ } ∈ Functional⊣]);

```

SML

```

a(PC_T1 "hol1"rewrite_tac[get_spec⊢ ColCons⊢,get_spec⊢ Num⊢]);
a(strip_asm_tac (list_∇_elim[⊢ clear⊢,⊢ z'⊢]cleanColCons_fun_thm));
a(POP_ASM_T rewrite_thm_tac);
a(∇_tac THEN rewrite_tac[cleanRows_def]);
a(DROP_NTH_ASM_T 7 ante_tac);
a(rewrite_tac[get_spec⊢ RowS⊢] THEN conv_tac(MAP_C let_conv)
  THEN PC_T1 "hol1"rewrite_tac[↔_def,×_def,get_spec⊢ Num⊢,get_spec⊢ DataS⊢,↔_def]);
a(lemma_tac⊢∃ l • TS_rows z' = l⊢ THEN LIST
  [prove_∃_tac,POP_ASM_T rewrite_thm_tac]);

```

SML

```

a(LIST_INDUCTION_T⊢l⊢asm_tac);
(* *** Goal "2.1" *** *)
a(rewrite_tac[⊢_def,map_def,map_null_thm,∈l_def]);
(* *** Goal "2.2" *** *)
a(∇_tac THEN rewrite_tac[map_def,⊢_def,∈l_def]);
a strip_tac;
a(lemma_tac⊢∇ r • r ∈l l ⇒ R_data r ∈ Functional⊢);
(* *** Goal "2.2.1" *** *)
a(REPEAT strip_tac);
a(cases_tac⊢r' = x''⊢);

```

SML

```

(* *** Goal "2.2.1.1" *** *)
a(DROP_NTH_ASM_T 3(ante_tac o  $\forall$ _elim $\lceil$ x'' $\rceil$ ) THEN asm_rewrite_tac[]);
(* *** Goal "2.2.1.2" *** *)
a(DROP_NTH_ASM_T 3(ante_tac o  $\forall$ _elim $\lceil$ r'' $\rceil$ ) THEN asm_rewrite_tac[]);
(* *** Goal "2.2.2" *** *)
a(cases_tac $\lceil$ clear dominates R_exist x'' $\rceil$ THEN asm_rewrite_tac[map_def, $\in$ _def]
  THEN strip_tac );
(* *** Goal "2.2.2.1" *** *)
a(POP_ASM_T rewrite_thm_tac THEN
  rewrite_tac[cleanRow_def,row_components,get_spec $\lceil$ MkRow $\rceil$ ]);
a(lemma_tac $\lceil$ filterRow (Snd (cleanColCons clear z')) (R_data x'')  $\in$  Functional $\rceil$ );
(* *** Goal "2.2.2.1.1" *** *)
a(DROP_NTH_ASM_T 3 (ante_tac o  $\forall$ _elim $\lceil$ x'' $\rceil$ ));
a(rewrite_tac[functional_def] THEN strip_tac);
a(strip_tac THEN rewrite_tac[filterRow_def,cleanColCons_def, $\triangleright$ _thm, $\triangleleft$ _thm,dom_def]);
a(REPEAT strip_tac);
a(DROP_NTH_ASM_T 9 (strip_asm_tac o list_ $\forall$ _elim $\lceil$ x''' $\rceil$ , $\lceil$ w $\rceil$ , $\lceil$ z'' $\rceil$ ));
(* *** Goal "2.2.2.1.2" *** *)
a(strip_asm_tac( $\forall$ _elim $\lceil$ clear $\rceil$ replaceData_fun_thm));
a(strip_asm_tac(list_ $\forall$ _elim $\lceil$ filterRow (Snd (cleanColCons clear z')) (R_data x'') $\rceil$ ,
 $\lceil$ replaceData clear $\rceil$ }_fun_thm));

```

SML

```

(* *** Goal "2.2.2.2" *** *)
a(DROP_NTH_ASM_T 5 ante_tac THEN asm_rewrite_tac[]);
(* *** Goal "2.2.2.3" *** *)
a(DROP_NTH_ASM_T 5 ante_tac THEN asm_rewrite_tac[]);
val hideR_lemma = save_pop_thm"hideR_lemma";

```

#### 4.4 hide Lemma

We now use the results of sections 4.1, 4.2 and 4.3 to prove that *hide* is a member of *SecureHide*.

SML

```

val secureHide_lemma = save_thm("secureHide_lemma",
  rewrite_rule[secureHideR_lemma,hideR_lemma]hideR_hide_lemma);

```

HOL output

```

secureHide_lemma =  $\vdash$  hide  $\in$  secureHide

```

---

## 5 CLOSING DOWN

The following ProofPower instruction restores the previous proof context.

SML

|*pop-pc*();

## 6 THE THEORY fef010

### 6.1 Parents

*fef009*

### 6.2 Children

*fef011 fef040*

### 6.3 Constants

**secureHideR**  $(Class \times StateR \rightarrow StateR) \mathbb{P}$

### 6.4 Definitions

**secureHideR**  $\vdash \forall h$

- $h \in secureHideR$
- $\Leftrightarrow (\forall c_1 c_2 s_1 s_2$
- $h(c_1, s_1) = h(c_1, s_2) \wedge c_1 \text{ dominates } c_2$
- $\Rightarrow h(c_2, s_1) = h(c_2, s_2))$

### 6.5 Theorems

**repState\_consistent**

**absState\_consistent**

$\vdash Consistent$

$(\lambda (repState', absState')$

- $(\forall a \bullet absState' (repState' a) = a)$
- $\wedge (\forall r$
- $isState r \Rightarrow repState' (absState' r) = r)$
- $\wedge (\forall a_1 a_2$
- $repState' a_1 = repState' a_2 \Leftrightarrow a_1 = a_2)$
- $\wedge (\forall r_1 r_2$
- $isState r_1 \wedge isState r_2$
- $\Rightarrow (absState' r_1 = absState' r_2$
- $\Leftrightarrow r_1 = r_2))$
- $\wedge (\forall s \bullet isState (repState' s)))$

**dir\_components**

$\vdash \forall dir_1 dir_2$

- $dir_1 = dir_2$
- $\Leftrightarrow Dir\_tables dir_1 = Dir\_tables dir_2$
- $\wedge Dir\_exist dir_1 = Dir\_exist dir_2$
- $\wedge Dir\_class dir_1 = Dir\_class dir_2$

**tab\_components**

$\vdash \forall t_1 t_2$

- $t_1 = t_2$



$$\begin{aligned}
&\Leftrightarrow TS\_class\ t_1 = TS\_class\ t_2 \\
&\wedge TS\_maxRow\ t_1 = TS\_maxRow\ t_2 \\
&\wedge TS\_colspecs\ t_1 = TS\_colspecs\ t_2 \\
&\wedge TS\_cons\ t_1 = TS\_cons\ t_2 \\
&\wedge TS\_rows\ t_1 = TS\_rows\ t_2
\end{aligned}$$

**row\_components**

$$\begin{aligned}
&\vdash \forall r_1\ r_2 \\
&\bullet r_1 = r_2 \\
&\Leftrightarrow R\_exist\ r_1 = R\_exist\ r_2 \\
&\wedge R\_data\ r_1 = R\_data\ r_2
\end{aligned}$$

**data\_components**

$$\begin{aligned}
&\vdash \forall d_1\ d_2 \\
&\bullet d_1 = d_2 \\
&\Leftrightarrow Dat\_class\ d_1 = Dat\_class\ d_2 \\
&\wedge Dat\_item\ d_1 = Dat\_item\ d_2
\end{aligned}$$

**hideR\_hide\_lemma**

$$\begin{aligned}
&\vdash hideR \in secureHideR \\
&\wedge (\forall clear\ s \\
&\bullet isState\ s \Rightarrow isState\ (hideR\ (clear, s))) \\
&\Rightarrow hide \in secureHide
\end{aligned}$$

**replaceData\_lemma**

$$\begin{aligned}
&\vdash \forall c_1\ c_2 \\
&\bullet c_1\ dominates\ c_2 \\
&\Rightarrow (\forall d_1\ d_2 \\
&\bullet replaceData\ c_1\ d_1 = replaceData\ c_1\ d_2 \\
&\Rightarrow replaceData\ c_2\ d_1 = replaceData\ c_2\ d_2)
\end{aligned}$$

**cleanColCons\_lemma**

$$\begin{aligned}
&\vdash \forall c_1\ c_2\ t_1\ t_2 \\
&\bullet c_1\ dominates\ c_2 \\
&\wedge cleanColCons\ c_1\ t_1 = cleanColCons\ c_1\ t_2 \\
&\Rightarrow cleanColCons\ c_2\ t_1 = cleanColCons\ c_2\ t_2
\end{aligned}$$

 **$\subseteq$ \_cleanColCons\_lemma**

$$\begin{aligned}
&\vdash \forall c_1\ c_2\ t \\
&\bullet c_1\ dominates\ c_2 \\
&\Rightarrow Fst\ (cleanColCons\ c_2\ t) \\
&\subseteq Fst\ (cleanColCons\ c_1\ t) \\
&\wedge Snd\ (cleanColCons\ c_2\ t) \\
&\subseteq Snd\ (cleanColCons\ c_1\ t)
\end{aligned}$$

 **$\subseteq$ \_cleanColCons\_lemma1**

$$\begin{aligned}
&\vdash \forall c_1\ c_2\ t\ col \\
&\bullet c_1\ dominates\ c_2 \\
&\Rightarrow col \in Snd\ (cleanColCons\ c_2\ t) \\
&\Rightarrow col \in Snd\ (cleanColCons\ c_1\ t)
\end{aligned}$$

**cleanRow\_lemma**

$$\begin{aligned}
&\vdash \forall c_1\ c_2 \\
&\bullet c_1\ dominates\ c_2 \\
&\Rightarrow (\forall t_1\ t_2\ r_1\ r_2 \\
&\bullet cleanColCons\ c_1\ t_1 = cleanColCons\ c_1\ t_2
\end{aligned}$$

$$\begin{aligned}
& \wedge \text{cleanColCons } c_2 \ t_1 \\
& = \text{cleanColCons } c_2 \ t_2 \\
& \wedge \text{cleanRow} \\
& \quad c_1 \\
& \quad (\text{Snd } (\text{cleanColCons } c_1 \ t_1)) \\
& \quad r_1 \\
& = \text{cleanRow} \\
& \quad c_1 \\
& \quad (\text{Snd } (\text{cleanColCons } c_1 \ t_2)) \\
& \quad r_2 \\
& \Rightarrow \text{cleanRow} \\
& \quad c_2 \\
& \quad (\text{Snd } (\text{cleanColCons } c_2 \ t_1)) \\
& \quad r_1 \\
& = \text{cleanRow} \\
& \quad c_2 \\
& \quad (\text{Snd } (\text{cleanColCons } c_2 \ t_2)) \\
& \quad r_2)
\end{aligned}$$

**cleanTable\_lemma**

$$\begin{aligned}
& \vdash \forall c_1 \ c_2 \\
& \bullet c_1 \text{ dominates } c_2 \\
& \Rightarrow (\forall t_1 \ t_2 \\
& \bullet \text{cleanTable } c_1 \ t_1 = \text{cleanTable } c_1 \ t_2 \\
& \Rightarrow \text{cleanTable } c_2 \ t_1 = \text{cleanTable } c_2 \ t_2)
\end{aligned}$$

**cleanDirectory\_lemma**

$$\begin{aligned}
& \vdash \forall c_1 \ c_2 \\
& \bullet c_1 \text{ dominates } c_2 \\
& \Rightarrow (\forall d_1 \ d_2 \\
& \bullet \text{cleanDirectory } c_1 \ d_1 = \text{cleanDirectory } c_1 \ d_2 \\
& \Rightarrow \text{cleanDirectory } c_2 \ d_1 \\
& = \text{cleanDirectory } c_2 \ d_2)
\end{aligned}$$

 **$\subseteq$ \_dir\_lemma**

$$\begin{aligned}
& \vdash \forall c_1 \ c_2 \ s \ x \ y \\
& \bullet c_1 \text{ dominates } c_2 \\
& \Rightarrow (x, y) \in s \triangleright \{\text{dir} \mid c_2 \text{ dominates } \text{Dir\_exist } \text{dir}\} \\
& \Rightarrow (x, y) \in s \triangleright \{\text{dir} \mid c_1 \text{ dominates } \text{Dir\_exist } \text{dir}\}
\end{aligned}$$

**dir\_class\_preserved\_lemma**

$$\begin{aligned}
& \vdash \forall c \ c_1 \ s \ s_1 \ x \ z \ z_1 \\
& \bullet \text{cleanDirectory } c_1 \ z = \text{cleanDirectory } c_1 \ z_1 \\
& \quad \wedge (x, z) \\
& \quad \in s \triangleright \{\text{dir} \mid c_1 \text{ dominates } \text{Dir\_exist } \text{dir}\} \\
& \quad \wedge (x, z_1) \\
& \quad \in s_1 \triangleright \{\text{dir} \mid c \text{ dominates } \text{Dir\_exist } \text{dir}\} \\
& \Rightarrow (x, z_1) \\
& \quad \in s_1 \triangleright \{\text{dir} \mid c_1 \text{ dominates } \text{Dir\_exist } \text{dir}\}
\end{aligned}$$

**secureHideR\_lemma**

$$\vdash \text{hideR} \in \text{secureHideR}$$

**cleanDir\_fun\_thm**

$$\vdash \forall \text{clear} \bullet \text{Graph } (\text{cleanDirectory } \text{clear}) \in \text{Functional}$$

**replaceData\_fun\_thm** $\vdash \forall \text{clear} \bullet \text{Graph} (\text{replaceData } \text{clear}) \in \text{Functional}$ **cleanColCons\_fun\_thm** $\vdash \forall \text{clear } \text{tab}$ 

- $\text{TS\_cons } \text{tab} \in \text{Functional}$

 $\Rightarrow \text{Fst} (\text{cleanColCons } \text{clear } \text{tab}) \in \text{Functional}$ **tables\_fun\_thm** $\vdash \forall \text{clear } \text{dir}$ 

- $\text{Dir\_tables } \text{dir} \in \text{Functional}$

 $\Rightarrow \text{Dir\_tables} (\text{cleanDirectory } \text{clear } \text{dir})$  $\in \text{Functional}$ **hideR\_lemma**  $\vdash \forall \text{clear } s \bullet \text{isState } s \Rightarrow \text{isState} (\text{hideR} (\text{clear}, s))$ **secureHide\_lemma** $\vdash \text{hide} \in \text{secureHide}$

## 7 INDEX

<i>cleanColCons_def</i> .....	5
<i>cleanColCons_fun_thm</i> .....	27
<i>cleanColCons_lemma</i> .....	13
<i>cleanDirectory_def</i> .....	5
<i>cleanDirectory_lemma</i> .....	24
<i>cleanDir_fun_thm</i> .....	26
<i>cleanRows_def</i> .....	5
<i>cleanRow_def</i> .....	5
<i>cleanRow_lemma</i> .....	15
<i>cleanTable_def</i> .....	5
<i>cleanTable_lemma</i> .....	22
<i>data_components</i> .....	8
<i>dir_class_preserved_lemma</i> .....	25
<i>dir_components</i> .....	6
<i>dominates_trans</i> .....	5
<i>fef010</i> .....	4
<i>filterRow_def</i> .....	5
<i>hideR_def</i> .....	5
<i>hideR_hide_lemma</i> .....	10
<i>hideR_lemma</i> .....	30
<i>hide_def</i> .....	5
<i>replaceData_def</i> .....	5
<i>replaceData_fun_thm</i> .....	27
<i>replaceData_lemma</i> .....	11
<i>repState_absState_def</i> .....	5
<i>row_components</i> .....	8
<i>secureHideR_lemma</i> .....	26
<i>secureHideR</i> .....	9
<i>secureHide_def</i> .....	5
<i>tables_fun_thm</i> .....	27
<i>tab_components</i> .....	7
$\subseteq$ <i>cleanColCons_lemma1</i> .....	14
$\subseteq$ <i>cleanColCons_lemma</i> .....	13
$\subseteq$ <i>dir_lemma</i> .....	24